

[This question paper contains 4 printed pages.]

Your Roll No.....

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Sr. No. of Question Paper : 3374

Unique Paper Code : 32341502

Name of the Paper : Theory of Computation

Name of the Course : **B.Sc. (Hons.) Computer Science (CBCS-LOCF)**

Semester : V

Duration : 3 Hours

Maximum Marks : 75

Instructions for Candidates

1. Write your Roll No. on the top immediately on receipt of this question paper.
2. Question No. 1 (Section A) is compulsory.
3. Attempt any 4 Questions from Nos. 2 to 7 (Section B).
4. Parts of a question must be answered together.
5. Consider $S = \{a, b\}$ for all the questions unless specified otherwise.

Section A

1. (a) Describe the language for the following regular expression: $((a+b)b)^*$. 2
- (b) Construct a regular expression for language L of all words having **b** at the second and third position. 2
- (c) Consider the following Context Free Grammar (CFG): 2

$S \rightarrow AbB$

$A \rightarrow aA \mid \Lambda$

$B \rightarrow aB \mid bB \mid \Lambda$

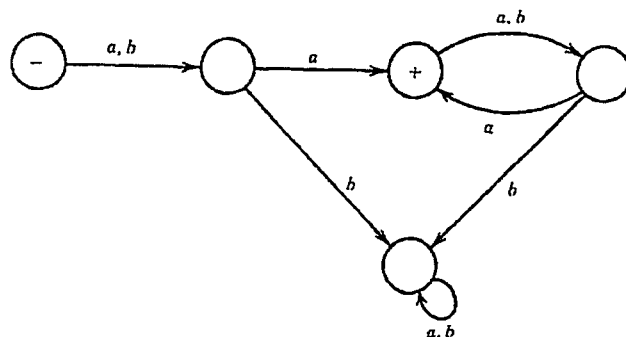
Create a parse tree for the word **aaba**

- (d) Give a context-free grammar (CFG) for b^*a^* . 3
- (e) Build a deterministic finite automaton that accepts all strings of **a**'s and **b**'s such that the next-to-last letter is an **a**. 3

P.T.O.

- (f) Describe the language accepted by the following finite automaton:

3



- (h) Using the pumping lemma, show that the following language is a non-regular language:

4

$$\{a^n b a^{2n} \text{ where } n \geq 1\} = \{abaa, aabaaaa, aaabaaaaaa, \dots\}$$

- (i) Build an FA that defines $L_1 \cap L_2$ where:

4

L_1 is the language of all odd-length words

$$L_2 = a(a+b)^*$$

- (j) Prove that a language is recursive if and only if the language itself and its complement are both recursively enumerable.

4

- (k) Show that the intersection of $L_1 = \{a^n b^m : n \geq m\}$ and $L_2 = \{a^n b^m : m \geq n\}$ is context-free.

4

- (l) Show that the following grammar is ambiguous:

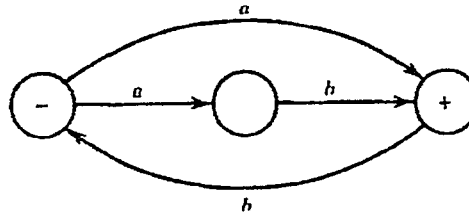
4

$$S \rightarrow XbaaX \mid aX$$

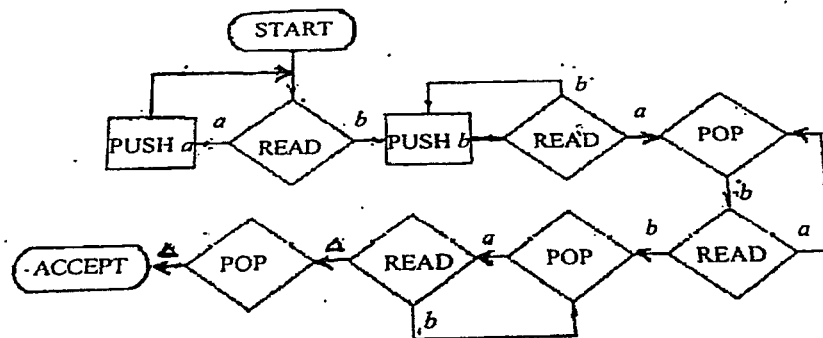
$$X \rightarrow Xa \mid Xb \mid \Lambda$$

Section B

- 2 (a) Convert the following non-deterministic finite automaton to a deterministic finite automaton: 5



- (b) i. Explain the halting problem. 5
 ii. Describe the Universal Turing Machine.
- 3 (a) Construct a PDA for the language $L = \{a^n b^n \mid n \geq 0\}$. 5
- (b) Is the language $\{a^n b^n c^m \mid n, m = 1, 2, 3, \dots\}$ context free? If yes, give a CFG 5
 for the same, else prove using the pumping lemma.
4. (a) Which language does the following PDA accept? Trace the string **aababb** on 6
 the PDA.



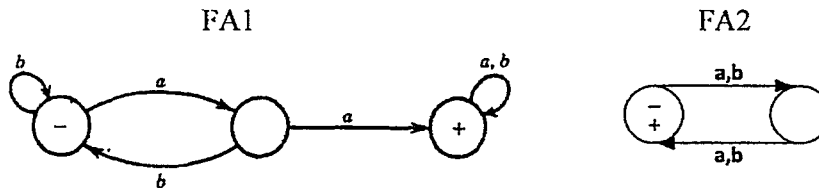
- (b) Convert the following CFG to Chomsky Normal Form: 4

$E \rightarrow E + E$
 $E \rightarrow E * E$
 $E \rightarrow D$
 $D \rightarrow 7$

The terminals are $+$, $*$, 7

- 5 (a) Find $FA1 \cap FA2$ corresponding to the following finite automata.

6



- (b) Construct a CFG for the language $L = a(bb)^*$. Also, construct a CFG for L^* . 4
- 6 (a) Assume $\Sigma = \{0, 1, \triangleright, \sqcup\}$, then create a Turing machine M_2 that takes a binary number as input on the tape, and leaves the two's complement of the input binary number as output on the tape when halts. 6
- (b) Are Context Free Languages closed under complementation? Justify. 4
- 7 (a) Build deterministic finite automata for the following languages over the alphabet $\{a, b\}$: 6
- (i) language of words having a double letter in them
 - (ii) language of words having an odd number of a 's
- (b) Convert the following TG into a regular expression: 4

